

1. 研究背景
2. 符号和定义
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Fractional matching preclusion

——2018年广东省组合图论会议

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1. Matching preclusion

- **A matching of G :** A function $f : E(G) \rightarrow \{0, 1\}$ such that for each vertex v , $\sum_{e \sim v} f(e) \leq 1$ where the sum is taken over all edges e incident with v , denoted by $e \sim v$.

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- **A perfect matching of G** : A matching satisfying that $\sum_{e \sim v} f(e) = 1$ for every vertex v of G .
- **An almost-perfect matching of G** : A matching satisfying that there is exactly one vertex v' such that $\sum_{e \sim v'} f(e) = 0$.

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$$\begin{aligned} \min & 1^T y \\ \text{s.t.} & y^T q^M \geq 1 \text{ for every } M \in \mathcal{M} \end{aligned} \quad (1.1)$$

$$y_e \in \{0, 1\} \text{ for every } e \in E(G) \quad (1.2)$$

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1. 研究背景

The concept of matching preclusion was introduced by Brigham[1]. As a measure of robustness in the event of edge failure in interconnection networks, as well as a theoretical connection to conditional connectivity, "changing and unchanging of invariants" and extremal graph theory. An interconnection network with a larger matching preclusion number may be considered as more robust in the event of link failures. The following results can be found in [1,5].

1.研究背景

Theorem 1[1]. Let $n \geq 2$ be an integer. Then

$$mp(K_n) = \begin{cases} n - 1 & n \text{ is even} \\ 2n - 3 & \text{otherwise} \end{cases}$$

Moreover, if $n \geq 6$ is even, then every optimal matching preclusion set is trivial. If $n \geq 11$ is odd, then every optimal matching preclusion set F is trivial and $G - F$ has two isolated vertices.

1. 研究背景

Theorem 2[5]. Let $n \geq 4$ be an even integer. Then

$$mp_1(K_n) = \begin{cases} \frac{(n^2+2n)}{8} & \text{if } n \in \{4, 6, 8\}, \\ 2n - 5, & \text{if } n \geq 10 \text{ and } n \text{ is even.} \end{cases}$$

Moreover, if $n \geq 12$, then every optimal conditional matching preclusion set F satisfies that $G - F$ has 2 leaves adjacent to the same vertex.

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Theorem 2[5]. Let $n \geq 5$ be odd integer. Then

$$mp_1(K_n) = \begin{cases} \frac{(n^2+4n+3)}{8}, & \text{if } n \in \{5, 7, 9, 11, 13\}, \\ 3n - 9, & \text{if } n \geq 15 \text{ and } n \text{ is odd.} \end{cases}$$

Moreover, if $n \geq 17$, then every optimal matching preclusion set F satisfies that $G - F$ has 3 leaves adjacent to the same vertex.

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- $smp(G)$:

$$smp(G) = \min\{|F| : F \text{ is an SMP set}\}.$$

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对于任意 $n \geq 8$ and $k \geq 3$ with $n - k \geq 5$

$$\text{smp}(S_{(n,k)}) = n - 1.$$

Moreover, every optimal strong matching preclusion set is trivial.

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- 在文献[8]中Ruizhi Lin 和Heping Zhang 修改了MP 模型, 得到下列FMP模型:

$$\begin{aligned} \min & \mathbf{1}^T \mathbf{y} \\ \text{s.t.} & \mathbf{y}^T \mathbf{q}^M \geq 1 \quad \text{for every } M \in \mathcal{M} \end{aligned} \quad (1.1)$$

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The optimal objective value of *FMP* is called the fractional matching preclusion number of G , denoted by $mp_f(G)$.

For any bipartite graph G , they give a formula for $mp_f(G)$ and study the relation between $mp_f(G)$ and $mp(G)$.

Moreover, they give a lower bound for $mp_f(G)$ of Cartesian product of two bipartite graphs

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3.1. 完全图的分數（強）匹配排除集和条件分數匹配排除数

- When $|V(G)|$ is even, $mp_1(G) \leq fmp_1(G) \leq 2n - 5$. So for a graph G with even vertices, if $mp_1(G) = 2n - 5$, then $mp_1(G) = fmp_1(G)$.

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- 参考文献[7]中给出：
- $fmp(K_n) = n - 1$ for every integer $n \geq 7$.
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- **定理1** $fmp(K_n) = n - 1$, for every even $n \geq 6$. Moreover, the optimal fractional matching preclusion set is trivial.

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- **定理1** $fmp(K_n) = n - 1$, for every even $n \geq 6$. Moreover, the optimal fractional matching preclusion set is trivial.
- **定理2** $fmp(K_n) = n - 1$, for every odd $n \geq 9$. Moreover, the optimal fractional matching preclusion set is trivial.

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3.1. 完全图的分數（強）匹配排除集和条件分數匹配排除数

- **定理3** $f_{smp}(K_n) = n - 2$, for every odd $n \geq 9$. Moreover, the optimal fractional strong matching preclusion set is $X_1 \cup \{e\}$ and $X_2 \cup E_3$ (Let $X_1 \subseteq V(K_n)$ such that $|X_1| = n - 1$, Then $K_n - X = K_3, e \in E(K_3)$. Let $X_2 \subseteq V(K_n)$ such that $|X_2| = n - 5$, Then $K_n - X = K_5, E_3$ is the edge set of a complete subgraph on 3 vertices).

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- **定理4** Let $n \geq 5$ be odd integer. Then

$$mp_1(K_n) = \begin{cases} \frac{(n^2-1)}{8}, & \text{if } n \in \{5, 7, 9, 11\}, \\ 2n - 5, & \text{if } n \geq 13 \text{ and } n \text{ is odd.} \end{cases} \quad (1)$$

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- **定理5** Let $n \geq 5$ be even integer. Then

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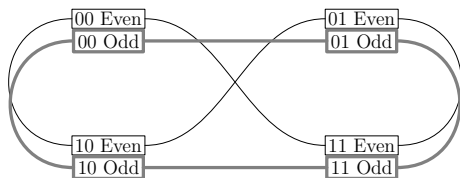
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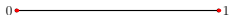
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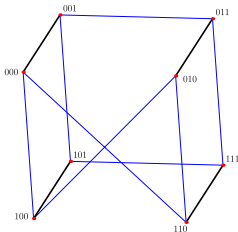
Cross edges

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1. 符号和定义



twisted 1-cube,



twisted 3-cube

3.3. 扭曲立方体的分数强匹配排除数

- **定理6** Let $n \geq 5$ be odd. Then $fsm_p(TQ_n) = n$ for every odd $n \geq 5$. Moreover, the optimal fractional strong matching preclusion set is trivial. .

3.4. (n,k) -star graph 的分数强匹配排除数和分数强匹配排除集

- The (n,k) -star graph, denoted $S_{(n,k)}$ is defined for positive integers n and k such that $n > k \geq 1$. The vertex set of the graph is all the permutations on k elements of the set $\{1, 2, \dots, n\}$. Two vertices corresponding to the permutations $[a_1, a_2, \dots, a_k]$ and $[b_1, b_2, \dots, b_k]$ are adjacent if and only if either:

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3.4. (n, k) -star graph 的分数强匹配排除数和分数强匹配排除集

- (1) There exists an integer $2 \leq s \leq k$ such that $a_1 = b_s$ and $b_1 = a_s$ and for any $i \neq s, 1 < i \leq k$, we have $a_i = b_i$. (That is, $[b_1, b_2, \dots, b_k]$ is obtained from $[a_1, a_2, \dots, a_k]$ by swapping a_1 and a_s .)

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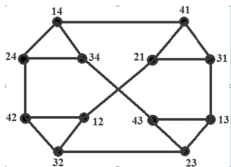
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- (2) For all $2 \leq i \leq k$, we have $a_i = b_i$ and $a_1 \neq b_1$. (That is, $[b_1, b_2, \dots, b_k]$ is obtained from $[a_1, a_2, \dots, a_k]$ by replacing a_1 by an element in $\{1, 2, \dots, n\} - \{a_1, a_2, \dots, a_k\}$.)

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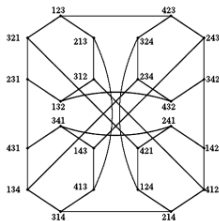
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1. 研究背景
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3.4. (n,k) -star graph 的分数强匹配排除数和分数强匹配排除集



$S(4,2)$,



$S(4,3)$

3.4. (n,k) -star graph 的分数强匹配排除数和分数强匹配排除集

- **定理7** $fsm_p(S_{(n,2)}) = n - 1$. Moreover every optimal fractional strong matching set is trivial, P_2 -trivial.

3.4. (n,k) -star graph 的分数强匹配排除数和分数强匹配排除集

- **定理7** $fsm_p(S_{(n,2)}) = n - 1$. Moreover every optimal fractional strong matching set is trivial, P_2 -trivial.
- **定理8** $fsm_p(S_{(n,3)}) = n - 1$. Moreover every optimal fractional strong matching set is trivial.

3.4. (n,k) -star graph 的分数强匹配排除数和分数强匹配排除集

- **定理7** $fsm_p(S_{(n,2)}) = n - 1$. Moreover every optimal fractional strong matching set is trivial, P_2 -trivial.
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- **定理9** Let $n > k \geq 3$ with $n - k \geq 2$, $fsm_p(S_{(n,k)}) = n - 1$. Moreover every optimal fractional strong matching set is trivial.

3.4. (n,k) -star graph 的分数强匹配排除数和分数强匹配排除集

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Thanks for your attention!